Quantum Information Summer School 2019, IBA

M. Sohaib Alam Rigetti Computing 15 July, 2019







The world's first **full-stack quantum computing** company.

16-qubit QPUs currently operating on our cloud platform

100+ employees w/ \$119M raised

Home of Fab-1, the world's first commercial quantum integrated circuit fab

Located in Berkeley, Calif. (R&D Lab) and Fremont, Calif.



Classical computers have fundamental limits



Transistor scaling

Economic limits with 10bn for next node fab

Ultimate single-atom limits



Returns to parallelization

Amdahl's law



Energy consumption

Exascale computing project has its own power plant

Power density can melt chips



Why build a quantum computer?

Quantum computing power* scales exponentially with qubits N bits can exactly simulate log N qubits

This compute unit....



Commodore 64



AWS M4 Instance

1 Million x Commodore 64



Entire Global Cloud

1 Billion x (1 Million x Commodore 64)

can exactly simulate:

10 Qubits

30 Qubits

60 Qubits



* More precisely ...

Why build a quantum computer?

For **N qubits** every time step (~100ns*) is an exponentially large **2^N x 2^N** complex **matrix multiplication**

Crucial details:

- limited number of multiplications (hundreds to thousands) due to noise
- not arbitrary matrices (need to be easily constructed on a QC)
- small I/O, poly(N)-bits in and N-bits out

The "big-memory small pipe" mental model for quantum computing





Quantum Advantage

Using quantum computers to solve problems **better**, faster, or cheaper than otherwise possible.

Marking a turning point in the way we do science and the way we do business.

Qubit:
$$|\psi\rangle = \alpha |0\rangle + \beta |1\rangle, \quad \alpha, \beta \in \mathbb{C}, \quad |\alpha|^2 + |\beta|^2 = 1$$
Kets: $|0\rangle = \begin{pmatrix} 1\\0 \end{pmatrix} \quad |1\rangle = \begin{pmatrix} 0\\1 \end{pmatrix} \quad |\psi\rangle = \begin{pmatrix} \alpha\\\beta \end{pmatrix}$

Measurement yields:

- 'O' with probability $|lpha|^2$
- '1' with probability $|eta|^2$



Qubit:
$$|\psi\rangle = \alpha |0\rangle + \beta |1\rangle, \quad \alpha, \beta \in \mathbb{C}, \quad |\alpha|^2 + |\beta|^2 = 1$$
Kets: $|0\rangle = \begin{pmatrix} 1\\0 \end{pmatrix} \quad |1\rangle = \begin{pmatrix} 0\\1 \end{pmatrix} \quad |\psi\rangle = \begin{pmatrix} \alpha\\\beta \end{pmatrix}$ Bras $\langle 0| = (1, 0) \quad \langle 1| = (0, 1) \quad \langle \psi| = (\overline{\alpha}, \overline{\beta})$

Brackets (Inner Product)

$$\begin{split} |\phi\rangle &= \gamma |0\rangle + \delta |1\rangle \\ \langle\phi|\psi\rangle &= \overline{\gamma}\alpha + \overline{\delta}\beta = \overline{\langle\psi|\phi\rangle} \end{split}$$



Qubit:
$$|\psi\rangle = \alpha |0\rangle + \beta |1\rangle, \quad \alpha, \beta \in \mathbb{C}, \quad |\alpha|^2 + |\beta|^2 = 1$$
Kets: $|0\rangle = \begin{pmatrix} 1\\0 \end{pmatrix} \quad |1\rangle = \begin{pmatrix} 0\\1 \end{pmatrix} \quad |\psi\rangle = \begin{pmatrix} \alpha\\\beta \end{pmatrix}$ Bras $\langle 0| = (1, 0) \quad \langle 1| = (0, 1) \quad \langle \psi| = (\overline{\alpha}, \overline{\beta})$

Normalization:
$$\langle \psi | \psi
angle = 1 \quad \Rightarrow \quad | \alpha |^2 + | \beta |^2 = 1$$



(

Qubit:
$$|\psi\rangle = \alpha |0\rangle + \beta |1\rangle,$$
 $\alpha, \beta \in \mathbb{C},$ $|\alpha|^2 + |\beta|^2 = 1$ Kets: $|0\rangle = \begin{pmatrix} 1\\0 \end{pmatrix}$ $|1\rangle = \begin{pmatrix} 0\\1 \end{pmatrix}$ $|\psi\rangle = \begin{pmatrix} \alpha\\\beta \end{pmatrix}$

Measurement yields:

- 'O' with probability $|\langle 0|\psi
 angle|^2=|lpha|^2$
- '1' with probability

 $|\langle 0|\psi\rangle|^2 = |\beta|^2$



Qubit States: Bloch Sphere



 $|\psi\rangle = \cos(\theta/2)|0\rangle + e^{i\phi}\sin(\theta/2)|1\rangle$



Qubit States: Multi-qubit states

Multiple qubits:

$$|\psi\rangle_{n-1}\otimes...\otimes|\psi\rangle_2\otimes|\psi\rangle_1\otimes|\psi\rangle_0$$

Tensor product:

$$\begin{aligned} |\psi\rangle \otimes |\phi\rangle &= (\alpha_0|0\rangle + \alpha_1|1\rangle) \otimes (\beta_0|0\rangle + \beta_1|1\rangle) \\ &= \alpha_0\beta_0|00\rangle + \alpha_0\beta_1|01\rangle + \alpha_1\beta_0|10\rangle + \alpha_1\beta_1|11\rangle \end{aligned}$$

Vector form:

$$\begin{pmatrix} \alpha_0 \\ \alpha_1 \end{pmatrix} \otimes \begin{pmatrix} \beta_0 \\ \beta_1 \end{pmatrix} = \begin{pmatrix} \alpha_0 \begin{pmatrix} \beta_0 \\ \beta_1 \\ \alpha_1 \begin{pmatrix} \beta_0 \\ \beta_1 \end{pmatrix} \end{pmatrix} = \begin{pmatrix} \alpha_0 \beta_0 \\ \alpha_0 \beta_1 \\ \alpha_1 \beta_0 \\ \alpha_1 \beta_1 \end{pmatrix}$$



Qubit States: Multi-qubit states

Associative:

$$(|\psi
angle\otimes|\phi
angle)\otimes|\gamma
angle=|\psi
angle\otimes(|\phi
angle\otimes|\gamma
angle)$$

Not commutative:

 $|\psi\rangle\otimes|\phi\rangle\neq|\phi\rangle\otimes|\psi\rangle$



Qubit Operations

Unitary Operators: (Gates)

$$UU^{\dagger}=U^{\dagger}U=I$$

$$|\psi\rangle \rightarrow |\psi'\rangle = U|\psi\rangle$$

Preserve inner product:

$$\langle \phi | \rightarrow \langle \phi' | = \langle \phi | U^{\dagger}$$

$$\langle \phi | \psi \rangle \rightarrow \langle \phi' | \psi' \rangle = \langle \phi | U^{\dagger} U | \psi \rangle = \langle \phi | \psi \rangle$$



	Bits	Probabilistic Bits		
State (single unit)	$Bit \in \{0,1\}$	Real vector $ec{b}=aec{0}+bec{1}$	$a+b \in \mathbb{R}_+$ $a+b=1$	

Qubits





Qubits

	Bits	Probabilisti	c Bits	Qubits	
State (single unit)	$Bit \in \{0, 1\}$	Real vector	$a+b \in \mathbb{R}_+$	Complex vector	$\alpha,\beta\in\mathbb{C}$
		$\vec{b} = a\vec{0} + b\vec{1}$	a + b = 1	$\vec{\psi} = \alpha \vec{0} + \beta \vec{1}$	$\left \alpha\right ^{2}+\left \beta\right ^{2}=1$





























	Bits	Probabilisti	c Bits	Qubits	
State (single unit)	Bit $\in \{0, 1\}$	Real vector	$a+b \in \mathbb{R}_+$	Complex vector	$\alpha,\beta\in\mathbb{C}$
		$\vec{b} = a\vec{0} + b\vec{1}$	a+b=1	$\vec{\psi} = \alpha \vec{0} + \beta \vec{1}$	$\left \alpha\right ^{2}+\left \beta\right ^{2}=1$
State (multi-unit)	Bitstring	Prob. Distribution (stochast	ic vector)		
	$x \in \{0,1\}^n$	$\vec{s} = \{p_x\}_{x \in \{0,1\}}$	n		
$\vec{s} = \bigotimes_{i}^{n} b_{i}$ Probability of bitstring x					

	Bits	Probabilistic Bits		Qubits	
State (single unit)	Bit $\in \{0, 1\}$	Real vector	$a+b \in \mathbb{R}_+$	Complex vector	$\alpha,\beta\in\mathbb{C}$
		$\vec{b} = a\vec{0} + b\vec{1}$	a+b=1	$\vec{\psi} = \alpha \vec{0} + \beta \vec{1}$	$\left \alpha\right ^{2}+\left \beta\right ^{2}=1$
State (multi-unit)	Bitstring	Prob. Distribution (stochastic vector)		Wavefunction (complex vec	ctor)
	$x \in \{0,1\}^n$	$\vec{s} = \{p_x\}_{x \in \{0,1\}^n}$		$\vec{\psi} = \{\alpha_x\}_{x \in \{0\}}$	$0,1\}^n$

$$\vec{s} = \bigotimes_{i}^{n} b_{i}$$

$$\vec{\psi} = \bigotimes_{i}^{n} \psi_{i}$$



	Bits	Probabilistic Bits		Qubits	
State (single unit)	Bit $\in \{0, 1\}$	Real vector	$a+b \in \mathbb{R}_+$	Complex vector	$\alpha,\beta\in\mathbb{C}$
		$\vec{b} = a\vec{0} + b\vec{1}$	a+b=1	$\vec{\psi} = \alpha \vec{0} + \beta \vec{1}$	$\left \alpha\right ^{2}+\left \beta\right ^{2}=1$
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		$\vec{s} = \bigotimes_{i}^{n} b_{i}$	α _x ² = Probabi	$\vec{\psi} = \bigotimes_{i}^{n}$ lity of bitstring x	ψ_i



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	$x \in \{0,1\}^n$	$\vec{s} = \{p_x\}_{x \in \{0,1\}^n}$		$\vec{\psi} = \{\alpha_x\}_{x \in \{}$	$[0,1]^n$
Operations	Boolean Logic	Stochastic Matrices			
		$\sum_{j=1}^S P_{i,j} = 1.$			





	Bits	Probabilistic Bits		Qubits	
State (single unit)	Bit $\in \{0, 1\}$	Real vector	$a+b\in\mathbb{R}_+$	Complex vector	$\alpha,\beta\in\mathbb{C}$
		$\vec{b} = a\vec{0} + b\vec{1}$	a+b=1	$\vec{\psi} = \alpha \vec{0} + \beta \vec{1}$	$\left \alpha\right ^{2}+\left \beta\right ^{2}=1$
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	$x \in \{0, 1\}^n$	$\vec{s} = \{p_x\}_{x \in \{0,1\}^n}$		$\vec{\psi} = \{\alpha_x\}_{x \in \{}$	$0,1\}^{n}$
Operations	Boolean Logic	Stochastic Matrices		Unitary Matrices	
		$\sum_{j=1}^S P_{i,j} = 1.$		$U^{\dagger}U =$	1







	Bits	Probabilistic Bits		Qubits	
State (single unit)	Bit $\in \{0, 1\}$	Real vector	$a+b\in\mathbb{R}_+$	Complex vector	$\alpha,\beta\in\mathbb{C}$
		$\vec{b} = a\vec{0} + b\vec{1}$	a+b=1	$\vec{\psi} = \alpha \vec{0} + \beta \vec{1}$	$\left \alpha\right ^{2}+\left \beta\right ^{2}=1$
State (multi-unit)	Bitstring	Prob. Distribution (stochastic	: vector)	Wavefunction (complex vector)	
	$x \in \{0,1\}^n$	$\vec{s} = \{p_x\}_{x \in \{0,1\}^n}$		$ec{\psi}=\{lpha_x\}_{x\in\{0,1\}^n}$	
Operations	Boolean Logic	Stochastic Matrices		Unitary Matrices	
		$\sum_{j=1}^S P_{i,j} = 1.$		$U^{\dagger}U =$	1
Component Ops	Boolean Gates	Tensor products of matrices		Tensor products of matrice	25



	Bits	Probabilistic Bits		Qubits			
State (single unit)	$Bit \in \{0, 1\}$	Real vector	$a+b \in \mathbb{R}_+$	Complex vector	$\alpha,\beta\in\mathbb{C}$		
		$\vec{b} = a\vec{0} + b\vec{1}$	a + b = 1	$\vec{\psi} = \alpha \vec{0} + \beta \vec{1}$	$\left \alpha\right ^{2}+\left \beta\right ^{2}=1$		
State (multi-unit)	Bitstring	Prob. Distribution (stochasti	c vector)	Wavefunction (complex vector)			
	$x \in \{0, 1\}^n$	$\vec{s} = \{p_x\}_{x \in \{0,1\}^n}$		$\vec{\psi} = \{\alpha_x\}_{x \in \mathbb{N}}$	$[0,1]^n$		
Operations	Boolean Logic	Stochastic Matrices		Unitary Matrices			
		$\sum_{j=1}^S P_{i,j} = 1.$		$U^{\dagger}U =$	1		
Component Ops	Boolean Gates	Tensor products of matrices	5	Tensor products of matric	es		
		Sampling					

	Bits	Probabilisti	: Bits	Qubits	
State (single unit)	$Bit \in \{0,1\}$	Real vector $ec{b}=aec{0}+bec{1}$	$a+b \in \mathbb{R}_+$ $a+b=1$	Complex vector $ec{\psi}=lphaec{0}+etaec{1}$	$lpha,eta\in\mathbb{C}$ $ lpha ^2+ eta ^2=1$
State (multi-unit)	Bitstring $x \in \{0,1\}^n$	Prob. Distribution (stochastic vector) $ec{s}=\{p_x\}_{x\in\{0,1\}^n}$		Wavefunction (complex vector) $ec{\psi}=\{lpha_x\}_{x\in\{0,1\}^n}$	
Operations	Boolean Logic	Stochastic Matrices $\sum_{j=1}^{S} P_{i,j} = 1.$		Unitary Matrices $U^{\dagger}U=$	1
Component Ops	Boolean Gates	Tensor products of matrice	S	Tensor products of matrice	25
Sampling $ \alpha_x ^2 = \text{Probability of bitstring x}$					

	Bits	Probabilistic Bits		Qubits	
State (single unit)	$Bit \in \{0, 1\}$	Real vector	$a+b \in \mathbb{R}_+$	Complex vector	$\alpha,\beta\in\mathbb{C}$
		$\vec{b} = a\vec{0} + b\vec{1}$	a+b=1	$\vec{\psi} = \alpha \vec{0} + \beta \vec{1}$	$\left \alpha\right ^{2}+\left \beta\right ^{2}=1$
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Operations	Boolean Logic	Stochastic Matrices		Unitary Matrices	
		$\sum_{j=1}^{S}P_{i,j}=1.$		$U^{\dagger}U =$	1
Component Ops	Boolean Gates	Tensor products of matrice	S	Tensor products of matric	es
	V=====	· · · · · · · · · · · · · · · · · · ·			
		Sampling	Bor	n rule	
				Measurem	nent

Qubit Operations: Pauli gates

$$I = \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} \qquad Y = \begin{pmatrix} 0 & -i \\ i & 0 \end{pmatrix}$$
$$X = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \qquad Z = \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix}$$



Qubit Operations: Identity gate

$$I|0\rangle = |0\rangle \qquad \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} \begin{pmatrix} 0 \\ 1 \end{pmatrix} = \begin{pmatrix} 0 \\ 1 \end{pmatrix}$$
$$I|1\rangle = |1\rangle \qquad \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} \begin{pmatrix} 1 \\ 0 \end{pmatrix} = \begin{pmatrix} 1 \\ 0 \end{pmatrix}$$



Qubit Operations: Pauli-X (NOT) gate

$$X|0\rangle = |1\rangle \qquad \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \begin{pmatrix} 1 \\ 0 \end{pmatrix} = \begin{pmatrix} 0 \\ 1 \end{pmatrix}$$
$$X|1\rangle = |0\rangle \qquad \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \begin{pmatrix} 0 \\ 1 \end{pmatrix} = \begin{pmatrix} 1 \\ 0 \end{pmatrix}$$



Qubit Operations: Pauli-Y gate

$$Y|0\rangle = i|1\rangle \qquad \begin{pmatrix} 0 & -i \\ i & 0 \end{pmatrix} \begin{pmatrix} 1 \\ 0 \end{pmatrix} = i \begin{pmatrix} 0 \\ 1 \end{pmatrix}$$
$$Y|1\rangle = -i|0\rangle \qquad \begin{pmatrix} 0 & -i \\ i & 0 \end{pmatrix} \begin{pmatrix} 0 \\ 1 \end{pmatrix} = -i \begin{pmatrix} 1 \\ 0 \end{pmatrix}$$



Qubit Operations: Pauli-Z gate

 $Z|0\rangle = |0\rangle \qquad \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix} \begin{pmatrix} 1 \\ 0 \end{pmatrix} = \begin{pmatrix} 1 \\ 0 \end{pmatrix}$ $Z|1\rangle = -|1\rangle \qquad \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix} \begin{pmatrix} 0 \\ 1 \end{pmatrix} = -\begin{pmatrix} 0 \\ 1 \end{pmatrix}$


Qubit Operations: Hadamard gate

$$H = \frac{1}{\sqrt{2}} \left(\begin{array}{cc} 1 & 1\\ 1 & -1 \end{array} \right)$$



Qubit Operations: Hadamard gate

$$H|0\rangle = \frac{1}{\sqrt{2}} \left(|0\rangle + |1\rangle\right) \qquad \frac{1}{\sqrt{2}} \left(\begin{array}{cc} 1 & 1\\ 1 & -1\end{array}\right) \left(\begin{array}{cc} 1\\ 0\end{array}\right) = \frac{1}{\sqrt{2}} \left(\begin{array}{cc} 1\\ 1\end{array}\right)$$
$$H|1\rangle = \frac{1}{\sqrt{2}} \left(|0\rangle - |1\rangle\right) \qquad \frac{1}{\sqrt{2}} \left(\begin{array}{cc} 1 & 1\\ 1 & -1\end{array}\right) \left(\begin{array}{cc} 0\\ 1\end{array}\right) = \frac{1}{\sqrt{2}} \left(\begin{array}{cc} 1\\ -1\end{array}\right)$$



Qubit Operations: Multiple qubits

Examples:

$$I \otimes X (|0\rangle \otimes |0\rangle) = |0\rangle \otimes |1\rangle = I_1 X_0 |00\rangle = |01\rangle$$

$$X \otimes H\left(|0\rangle \otimes |0\rangle\right) = |1\rangle \otimes \frac{1}{\sqrt{2}}\left(|0\rangle + |1\rangle\right) = X_1 H_0 |00\rangle = \frac{1}{2}\left(|10\rangle + |11\rangle\right)$$



Qubit Operations: Multiple qubits

$$A \otimes B = \begin{pmatrix} A_{00} & A_{01} \\ A_{10} & A_{11} \end{pmatrix} \otimes \begin{pmatrix} B_{00} & B_{01} \\ B_{10} & B_{11} \end{pmatrix} = \begin{pmatrix} A_{00}[B] & A_{01}[B] \\ A_{10}[B] & A_{11}[B] \end{pmatrix}$$
$$= \begin{pmatrix} A_{00}B_{00} & A_{00}B_{01} & A_{01}B_{00} & A_{01}B_{01} \\ A_{00}B_{10} & A_{00}B_{11} & A_{01}B_{10} & A_{01}B_{11} \\ A_{10}B_{00} & A_{01}B_{01} & A_{11}B_{00} & A_{11}B_{01} \\ A_{10}B_{10} & A_{11}B_{11} & A_{11}B_{10} & A_{11}B_{11} \end{pmatrix}$$



Qubit Operations: Quantum Circuits











from pyquil import Program, get_qc from pyquil.gates import X

p = Program(X(0))
qc = get_qc('9q-generic-qvm')
results = qc.run_and_measure(p, trials=10)[0]

print (results)

[111111111]





> [1] [1]]

from pyquil import Program, get_qc from pyquil.gates import X, MEASURE

<pre>p = Program() p.declare('ro', 'BIT', 1) p.inst(X(0)) p.inst(MEASURE(0, 'ro')) p.wrap_in_numshots_loop(shots=10)</pre>	
qc = get_qc(<mark>'9q-generic-qvm</mark> ') results = qc.run(qc.compile(p))	

print (results)



 $X|0\rangle = |1\rangle$



from pyquil import Program, get_qc from pyquil.gates import X, MEASURE

p = Program()
p.declare('ro', 'BIT', 1)
p.inst(X(0))
p.inst(MEASURE(0, 'ro'))
p.wrap_in_numshots_loop(shots=10)

qc = get_qc('9q-generic-qvm')
results = qc.run(qc.compile(p))

DECLARE ro BIT[1] X 0 MEASURE 0 ro[0]



print (p)

 $X_1H_0 |0\rangle_1|0\rangle_0$

from pyquil.gates import X, H, MEASURE

from pyquil import Program, get_qc



p = Program()[[0 1] ro = p.declare('ro', 'BIT', 2) [11]p.inst(H(0)) $[0\,1]$ p.inst(X(1))p.inst(MEASURE(0, ro[0])) [11]p.inst(MEASURE(1, ro[1])) [11]p.wrap in numshots loop(shots=10) $[0\,1]$ [11] $qc = get_qc('9q-generic-qvm')$ $[0\,1]$ results = qc.run(qc.compile(p)) $[0\,1]$ print (results) [0 1]]



 $Y_1 Z_0 X_1 H_0 |0\rangle_1 |0\rangle_0$



from pyquil import Program, get_qc
from pyquil.gates import X, Y, Z, H, MEASURE

[[10] [10]p = Program()ro = p.declare('ro', 'BIT', 2)[00] p += Program(H(0), X(1), Z(0), Y(1), MEASURE(0, ro[0]), MEASURE(1, ro[1])) [10]p.wrap_in_numshots_loop(shots=10) [10][00] qc = get_qc('9q-generic-qvm') [10]results = qc.run(qc.compile(p)) [10]print (results) [00] [[00]]



Quantum Dice

Goal: Create an N-sided dice using a quantum computer.



Quantum Dice

Question: What gate would we use?

$$H^{\otimes n}|0\rangle = \frac{1}{2^n} \sum_{z=0}^{n-1} |z\rangle$$



Quantum Dice

Question: How many qubits would we use?

$$2^n = N \Rightarrow n = \log_2 N$$



Qubit Operations: Projections

Project a ket/bra along a given ket/bra via its corresponding projection operator.

For some $\ket{\psi}$ the corresponding projection operator is given by the **outer product**

$$P_{\psi} = |\psi\rangle\langle\psi| = \begin{pmatrix} \alpha \\ \beta \end{pmatrix} (\overline{\alpha}, \ \overline{\beta}) = \begin{pmatrix} |\alpha|^2 & \alpha\overline{\beta} \\ \beta \overline{\alpha} & |\beta|^2 \end{pmatrix}$$

e.g.

$$P_{0} = \begin{pmatrix} 1 & 0 \\ 0 & 0 \end{pmatrix}, P_{1} = \begin{pmatrix} 0 & 0 \\ 0 & 1 \end{pmatrix}, P_{+} = \frac{1}{2} \begin{pmatrix} 1 & 1 \\ 1 & 1 \end{pmatrix}$$



Qubit Operations: Controlled Operations

$$c U = |0\rangle \langle 0| \otimes I + |1\rangle \langle 1| \otimes U$$

If qubit 1 is in the state |0>, apply *I* (identity) to qubit 0 Else if qubit 1 is in the state |1>, apply *U* to qubit 0

For example,

$$CNOT = |0\rangle\langle 0| \otimes I + |1\rangle\langle 1| \otimes X = \begin{pmatrix} 1 & 0 & 0 & 0\\ 0 & 1 & 0 & 0\\ 0 & 0 & 0 & 1\\ 0 & 0 & 1 & 0 \end{pmatrix}$$



Entangled States

A state that cannot be written as a product state, i.e.

$$|\psi\rangle \neq |\xi\rangle \otimes |\phi\rangle$$

An example of a state that is not entangled:

$$\frac{1}{\sqrt{2}}\left(|00\rangle + |01\rangle\right) = |0\rangle \otimes \frac{1}{\sqrt{2}}\left(|0\rangle + |1\rangle\right)$$

An example of a state that is entangled:

$$\frac{1}{\sqrt{2}}\left(\left|00\right\rangle+\left|11\right\rangle\right)$$



Bell State via CNOT



from pyquil import Program from pyquil.gates import H, CNOT from pyquil.api import WavefunctionSimulator

p = Program(H(1))
p += Program(CNOT(1, 0))
wfn = WavefunctionSimulator().wavefunction(p)

print (wfn)

(0.7071067812+0j)|00> + (0.7071067812+0j)|11>



Measurement in some arbitrary basis

Computational basis: $\{|0\rangle, |1\rangle\}$

Measurement yields:

- 'O' with probability $\langle \psi | P_0 | \psi \rangle = \langle \psi | 0 \rangle \langle 0 | \psi \rangle = |\langle 0 | \psi \rangle|^2$
- '1' with probability $\langle \psi|P_1|\psi
 angle=\langle\psi|1
 angle\langle 1|\psi
 angle=|\langle 1|\psi
 angle|^2$



Measurement in some arbitrary basis

Some other basis: $\{|0'
angle=U|0
angle, |1'
angle=U|1
angle\}$

Measurement yields:

- 0' with probability $\langle \psi | P_{0'} | \psi \rangle = \langle \psi | 0' \rangle \langle 0' | \psi \rangle = \langle \psi | U | 0 \rangle \langle 0 | U^{\dagger} | \psi \rangle$
- 1' with probability

$$= \langle \psi' | 0 \rangle \langle 0 | \psi' \rangle = |\langle 0 | \psi' \rangle|^{2}$$
$$\langle \psi | P_{1'} | \psi \rangle = \langle \psi | 1' \rangle \langle 1' | \psi \rangle = \langle \psi | U | 1 \rangle \langle 1 | U^{\dagger} | \psi \rangle$$
$$= \langle \psi' | 1 \rangle \langle 1 | \psi' \rangle = |\langle 1 | \psi' \rangle|^{2}$$

 $|\psi'\rangle = U^{\dagger}|\psi\rangle$



Measurement in some arbitrary basis

Measurement of $|\Psi\rangle$ in some basis {U|0>, U|1>}

= Measurement of U[†] $|\Psi\rangle$ in standard/computational basis { $|0\rangle$, $|1\rangle$ }



Goal: Teleport a Qubit!



Scenario:

- Alice is in possession of a qubit $|\Psi\rangle$, which she would like to teleport over to Bob, who is at some distant location.



Protocol:

- Create a Bell state, giving one qubit each to Alice and Bob
- Have Alice measure both her qubits in the Bell basis, and send her results to Bob
- Have Bob *conditionally* apply gates to his qubits, based off Alice's measurements, to reconstruct the original qubit at his location









Classical Control in pyQuil

```
from pyquil import Program
from pyquil.gates import I, X
from pyquil.api import WavefunctionSimulator
```

```
p = Program(X(0))
ro = p.declare('ro', 'BIT', 1)
p.measure(0, ro[0]).if_then(ro[0], Program(X(1)), Program(I(1)))
wfn = WavefunctionSimulator().wavefunction(p)
```

print (wfn)

(1+0j)|11>



Classical Control in pyQuil

```
from pyquil import Program
from pyquil.gates import I, X
from pyquil.api import WavefunctionSimulator
```

```
p = Program(I(0))
ro = p.declare('ro', 'BIT', 1)
p.measure(0, ro[0]).if_then(ro[0], Program(X(1)), Program(I(1)))
wfn = WavefunctionSimulator().wavefunction(p)
```

print (wfn)

(1+0j)|00>



Goal: Given a function f: $\{0, 1\} \rightarrow \{0, 1\}$, determine whether it is constant (i.e. f(0) = f(1)) or balanced (i.e. f(0) != f(1)) in the minimum number of steps, assuming an oracle/black-box for the function.



$$U_f|x\rangle|y\rangle = |x\rangle|y \oplus f(x)\rangle$$





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$$U_f\left(\frac{|0\rangle+|1\rangle}{\sqrt{2}}\right)\left(\frac{|0\rangle-|1\rangle}{\sqrt{2}}\right) = (-1)^{f(0)}\left(\frac{|0\rangle+(-1)^{f(0)\oplus f(1)}|1\rangle}{\sqrt{2}}\right)\left(\frac{|0\rangle-|1\rangle}{\sqrt{2}}\right)$$



(a)
$$f(0) = f(1)$$

$$U_f\left(\frac{|0\rangle + |1\rangle}{\sqrt{2}}\right)\left(\frac{|0\rangle - |1\rangle}{\sqrt{2}}\right) = \pm \left(\frac{|0\rangle + |1\rangle}{\sqrt{2}}\right)\left(\frac{|0\rangle - |1\rangle}{\sqrt{2}}\right)$$

(b) f(0) != f(1)

$$U_f\left(\frac{|0\rangle+|1\rangle}{\sqrt{2}}\right)\left(\frac{|0\rangle-|1\rangle}{\sqrt{2}}\right) = \pm\left(\frac{|0\rangle-|1\rangle}{\sqrt{2}}\right)\left(\frac{|0\rangle-|1\rangle}{\sqrt{2}}\right)$$

 $H|0\rangle = \left(\frac{|0\rangle + |1\rangle}{\sqrt{2}}\right), \quad H|1\rangle = \left(\frac{|0\rangle - |1\rangle}{\sqrt{2}}\right), \quad H^{\dagger} = H$

Recall:

Goal: Given a function f: $\{0, 1\}^n \rightarrow \{0, 1\}$, find the bitstring $x \in \{0, 1\}^n$, such that f(x) = 1.



Assume a quantum black box

$$U_f|x\rangle|y\rangle = |x\rangle|y \oplus f(x)\rangle$$

Prepare the query register as

$$\frac{1}{\sqrt{N}}\sum_{x=0}^{N-1}|x\rangle = \frac{1}{\sqrt{N}}|\psi_{good}\rangle + \sqrt{\frac{N-1}{N}}|\psi_{bad}\rangle$$

where

$$|\psi_{good}
angle=|w
angle$$
 ('w' is the bitstring to be found) $|\psi_{bad}
angle=rac{1}{\sqrt{N-1}}\sum_{x\in X_{bad}}|x
angle$



Oracle/Blackbox

$$\frac{\frac{1}{\sqrt{2^n}}|w\rangle + \sqrt{\frac{2^n-1}{2^n}}|\psi_{\text{bad}}\rangle = U_f = \int \frac{1}{\sqrt{2^n}}|w\rangle|1\rangle + \sqrt{\frac{2^n-1}{2^n}}|\psi_{\text{bad}}\rangle|0\rangle$$



Define a phase-shift operator

$$U_{\psi^{\perp}}: \left\{ \begin{array}{l} |x\rangle \to -|x\rangle, \quad \langle x|\psi\rangle = 0\\ |\psi\rangle \to |\psi\rangle \end{array} \right.$$

When $|\Psi\rangle$ is the equal-superposition state ('quantum dice'), we can write

$$U_{\psi^{\perp}} = 2|\psi\rangle\langle\psi| - I$$


Inversion about the mean:

$$U_{\psi^{\perp}} \sum_{i=0}^{2^{n}-1} \alpha_{i} |i\rangle = \sum_{i=0}^{2^{n}-1} (2\langle \alpha \rangle - \alpha_{i}) |i\rangle$$



Prepare target qubit as $ZH|0> = HX|0> \sim (|0> - |1>)$ to get phase 'kick back' from applying oracle/blackbox

$$U_f|x\rangle\left(\frac{|0\rangle-|1\rangle}{\sqrt{2}}\right) = (-1)^{f(x)}|x\rangle\left(\frac{|0\rangle-|1\rangle}{2}\right)$$



- Prepare the initial state

$$H^{\otimes n} \otimes (H_0 Z_0) | 00 \dots 0 \rangle$$

- Apply the Grover iterate $O(\sqrt{N})$ times

$$G = U_{\psi^{\perp}} U_f$$















Variational Quantum Eigensolver

Goal: Find the ground state of some Hamiltonian, H.

Procedure:

- Prepare some trial state $|\Psi; \theta >$
- Calculate expectation value $\langle \Psi; \theta | H | \Psi; \theta \rangle$
- Find the values of θ that will minimize the above

Goal: Given binary constraints over bitstrings

 $z \in \{0,1\}^n$

$$C_{\alpha}(z) = \begin{cases} 1 & \text{if } z \text{ satisfies the constraint } \alpha \\ 0 & \text{if } z \text{ does not} \end{cases}$$

Find the bitstring that maximizes the objective function

$$\operatorname{argmax}_{z} C(z) = \operatorname{argmax}_{z} \sum_{\alpha=1}^{m} C_{\alpha}(z)$$



MaxCut problem:

Given some undirected graph with arbitrary (non-negative) weights, find a partition (S, \overline{S}) of the graph's nodes (a 'cut' of the graph) that maximizes the weights along the cut

















- Prepare equal-superposition as initial state

$$|s\rangle = \frac{1}{\sqrt{2^n}} \sum_{z} |z\rangle$$

- Define the 'mixer' operator $B = \sum_{j=1}^{n} X_j$, $U(B, \beta) = e^{-i\beta B}$ - Define the 'cost' operator $C = \frac{1}{2} \sum_{i,j=1}^{n} w_{ij} \frac{1 - Z_i Z_j}{2}$, $U(C, \gamma) = e^{-i\gamma C}$
- Apply $U(B,\beta_p)U(C,\gamma_p)\dots U(B,\beta_1)U(C,\gamma_1)|s\rangle$ and sample/measure



import numpy as np from pyquil import Program from pyquil.api import WavefunctionSimulator from pyquil.paulis import sI, sX, sY, sZ, exponential_map

```
angle = np.pi / 8
pauli_sum = sX(1) * sY(0) + sI(1) * sY(0)
```

```
p = Program()
for ps in pauli_sum:
    p += exponential_map(ps)(angle)
```

```
wfn_sim = WavefunctionSimulator()
wfn = wfn_sim.wavefunction(p)
```

print (wfn)



'Pure' quantum states

$$|\psi\rangle = \alpha|0\rangle + \beta|1\rangle$$

evolve via Unitary operations

$$|\psi\rangle \rightarrow |\psi'\rangle = U|\psi\rangle \qquad UU^{\dagger} = U^{\dagger}U = I$$



More generally, quantum states are described by "Density Matrix"

$$\rho = \sum_{i} p_i |\psi_i\rangle \langle \psi_i$$

evolving via Kraus operations ("quantum channel")

$$\rho \to \sum_{i} K_{i} \rho K_{i}^{\dagger}, \qquad \sum_{i} K_{i}^{\dagger} K_{i} = I$$



For example,

$$\rho = \frac{1}{2} \left(|0\rangle \langle 0| + |1\rangle \langle 1| \right) = \frac{1}{2} \left(\begin{array}{cc} 1 & 0 \\ 0 & 1 \end{array} \right)$$

Not to be confused with

$$\rho = \left(\frac{1}{\sqrt{2}}(|0\rangle + |1\rangle)\right) \left(\frac{1}{\sqrt{2}}(\langle 0| + \langle 1|)\right) = \frac{1}{2} \left(\begin{array}{cc} 1 & 1\\ 1 & 1 \end{array}\right)$$



Example of quantum channel/set of Kraus operators/noise model:

$$\{\sqrt{p}X, \sqrt{1-p}Z\}$$

Quantum state passing through the channel/experiencing the noise transforms to:

$$\rho \to p X^{\dagger} \rho X + (1-p) Z^{\dagger} \rho Z$$



Thanks, and keep in touch!

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